



# The Future of Reverse Engineering with Large Language Models

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# About Us

- Tim Blazytko
  - Chief Scientist & Head of Engineering, co-founder of emproof
  - designs software protections for embedded devices
  - trainer for (de)obfuscation and reverse engineering techniques
- Moritz Schloegel
  - fresh postdoc at CISPA Helmholtz Center
  - working with bugs by day (mostly fuzzing)
  - code deobfuscation by night



## Setting the Scene

- ?
- Using LLMs for RE
- ✎ Local LLMs
- 🏆 Enhancements through Static Analysis

# LLMs in Reverse Engineering

## Disclaimer

- **hyped** and fast-developing field
- **teasing** powers and limitations for RE
- **not specific** to tools or LLMs

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- teasing powers and limitations for RE  
current snapshot, maybe soon outdated
- not specific to tools or LLMs

# Applications to Reverse Engineering

- renaming functions

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- renaming functions
- renaming variables

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- renaming functions
- renaming variables
- commenting code

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- renaming functions
- renaming variables
- commenting code
- explaining code

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- answering questions
- scripting support

# Applications to Reverse Engineering

- renaming functions
- renaming variables
- commenting code
- explaining today: focus on use cases
- answering questions
- scripting support

# Use Case Function Preselection

“For the given decompiler output, analyze the code and suggest a meaningful function name.”

Renamed function at 0x10002b50 to CallWithArguments

Renamed function at 0x10004b50 to DecodeComplexAlgorithm

Renamed function at 0x10008b60 to SetValueToMemoryLocation

Renamed function at 0x10002b70 to CallFunctionPointerWithArguments

Renamed function at 0x10002b50 to CallWithArguments  
Renamed function at 0x10004b50 to DecodeComplexAlgorithm  
Renamed function at 0x10006b50 to EncodeComplexAlgorithm  
Renamed function at 0x10008b50 to CallWithArguments

often too generic

Renamed function at 0x40cbb5 to `DecompileCodeAnalyze`  
Renamed function at 0x4033c2 to `DecompileCodeAnalyze`  
Renamed function at 0x4024c2 to `DecompileCodeAnalyze`  
Renamed function at 0x402d58 to `DecompileCodeAnalyze`  
Renamed function at 0x40ed62 to `DecompileAndProtectMemoryPage`  
Renamed function at 0x409e53 to `DecompileAndFindMatchingStringInMemory`  
Renamed function at 0x40e4dd to `DecompileAndAnalyzeFunction`

Renamed function at 0x40cbb5 to `DecompileCodeAnalyze`

Renamed function at 0x4033c2 to `DecompileCodeAnalyze`

Renamed function at 0x4024c2 to `DecompileCodeAnalyze`

Renamed function at 0x4024d2 to `DecompileCodeAnalyze`

Renamed function at 0x4024e2 to `DecompileCodeAnalyze`

sometimes entirely useless

Renamed function at 0x409e53 to `DecompileAndFindMatchingStringInMemory`

Renamed function at 0x40e4dd to `DecompileAndAnalyzeFunction`

Renamed function at 0x1000cf10 to `InitializeKeyLoggerAndHandleErrors`

Renamed function at 0x10014ae0 to `CreateNamedPipesAndRunShellCommands`

Renamed function at 0x100186d0 to `SearchForAProcessByName`

Renamed function at 0x1001d880 to `SendHTTPPOSTRequestAndHandleResponse`

Renamed function at 0x1000cf10 to InitializeKeyLoggerAndHandleErrors

Renamed function at 0x10014ae0 to CreateNamedPipesAndRunShellCommands

Renamed function at can be very helpfulame

Renamed function at andHandleResponse

# Use Case Identification of Library Functions

```
000007ac void* copy_backward_string(void* arg1, void* arg2, int32_t arg3)

000007ac     int32_t i = arg3
000007b0     void* r4 = arg2 - 1
000007b4     void* r3 = arg1 - 1
000007c0     do {
000007b8         char r5 = *(r4 + 1)
000007b8         r4 = r4 + 1
000007bc         *(r3 + 1) = r5
000007bc         r3 = r3 + 1
000007c0         i = i - 1
000007c0     } while (i != 0)
000007c4     return r3
```

Log    Q Search log

```
[Default] Renaming sub_7ac to copy_backward_string
[Default] Renaming sub_820 to return_address_of_data_838
[Default] Renaming sub_3e4 to process_input_and_copy_string
[Default] Renaming sub_77c to save_and_clear_fpu_registers
```

```
000007ac void* copy_backward_string(void* arg1, void* arg2, int32_t arg3)

000007ac     int32_t i = arg3
000007b0     void* r4 = arg2 - 1
000007b4     void* r3 = arg1 - 1
000007c0     do {
000007b8         char r5 = *(r4 + 1)
000007b8         r4 = r4 + 1
000007bc         *(r3 + 1) = r5
```

good approximation, but incorrect

```
000007c4     return r3
```

Log    Q Search log

```
[Default] Renaming sub_7ac to copy_backward_string
[Default] Renaming sub_820 to return_address_of_data_838
[Default] Renaming sub_3e4 to process_input_and_copy_string
[Default] Renaming sub_77c to save_and_clear_fpu_registers
```

# Enhancing Decompiler Output

“Enhance the following decompiler output by suggesting more meaningful variable names. Also, add comments.”

# memcpy

```
c                                         Copy

000007ac void* copy_memory(void* destination, void* source, int32_t count)
000007ac     int32_t remaining_bytes = count;           // Initialize remaining bytes to be copied
000007b0     void* source_ptr = source - 1;           // Source pointer starts before the actual source
000007b4     void* dest_ptr = destination - 1;         // Destination pointer starts before the actual
destination
000007c0     do {
000007b8         char byte_to_copy = *(source_ptr + 1); // Fetch the next byte from the source
000007b8         source_ptr = source_ptr + 1;           // Move the source pointer to the next byte
000007bc         *(dest_ptr + 1) = byte_to_copy;        // Write the byte to the destination
000007bc         dest_ptr = dest_ptr + 1;               // Move the destination pointer to the next byte
000007c0         remaining_bytes = remaining_bytes - 1; // Decrement the remaining byte count
000007c0     } while (remaining_bytes != 0);          // Repeat until all bytes are copied
000007c4     return dest_ptr;                         // Return the pointer to the last copied byte in
the destination
```

# memcpy

```
c
000007ac void* copy_memory(void* destination, void* source, int32_t count)
000007ac     int32_t remaining_bytes = count;           // Initialize remaining bytes to be copied
000007b0     void* source_ptr = source - 1;          // Source pointer starts before the actual source
000007b4     void* dest_ptr = destination - 1;        // Destination pointer starts before the actual
destination
000007c0
000007b8
000007b8     source_ptr = source_ptr + 1;           // Move the source pointer to the next byte
000007bc     *(dest_ptr + 1) = byte_to_copy;         // Write the byte to the destination
000007bc     dest_ptr = dest_ptr + 1;                 // Move the destination pointer to the next byte
000007c0     remaining_bytes = remaining_bytes - 1; // Decrement the remaining byte count
000007c0 } while (remaining_bytes != 0);           // Repeat until all bytes are copied
000007c4     return dest_ptr;                      // Return the pointer to the last copied byte in
the destination
```

works well for common code

Copy

# RC4 KSA

```
// Function to perform a transformation on input string based on some algorithm
char transform_string(char* input_string, int32_t param, int32_t* lookup_table) {
    char* source = input_string;
    char* index_ptr = nullptr; // Pointer used to iterate over the input string
    int32_t intermediate_values[1024] = {0}; // Buffer to store intermediate values
    void* buffer; // Temporary buffer to use during memset
    __builtin_memset(&buffer, 0, 0x3fc); // Initialize buffer with zeros
    int32_t accumulator = 0; // Accumulator to store intermediate sums
    int32_t* intermediate_ptr = intermediate_values;

    do {
        uint32_t temp_val = lookup_func(0:index_ptr, lookup_table); // Get value from lookup table
        *(index_ptr + source) = index_ptr->byte_val; // Assign byte value to source
        index_ptr = &index_ptr[1]; // Increment index pointer
        *intermediate_ptr = zero_extend(*(temp_val + param)); // Extend and store value in intermediate
array
        intermediate_ptr = &intermediate_ptr[1]; // Increment intermediate array pointer
    } while (index_ptr < 0x100); // Loop until index_ptr reaches 0x100
```

# RC4 KSA

```
// Function to perform a transformation on input string based on some algorithm
char transform_string(char* input_string, int32_t param, int32_t* lookup_table) {
    char* source = input_string;
    char* index_ptr = nullptr; // Pointer used to iterate over the input string
    int32_t intermediate_values[1024] = {0}; // Buffer to store intermediate values
    void* buffer; // Temporary buffer to use during memset
    __builtin_memset(&buffer, 0, 0x3fc); // Initialize buffer with zeros
    int32_t accumulator = 0;
    int32_t* intermediate_ptr = not always helpful
    do {
        uint32_t temp_val = lookup_func(0:index_ptr, lookup_table); // Get value from lookup table
        *(index_ptr + source) = index_ptr->byte_val; // Assign byte value to source
        index_ptr = &index_ptr[1]; // Increment index pointer
        *intermediate_ptr = zero_extend(*(temp_val + param)); // Extend and store value in intermediate
array
        intermediate_ptr = &intermediate_ptr[1]; // Increment intermediate array pointer
    } while (index_ptr < 0x100); // Loop until index_ptr reaches 0x100
```

# Use Case Analysis Guidance for Unfamiliar Instruction Sets

“Explain me the following assembly snippet.”

```
; Load Immediate Shifted:  
; Load the immediate value 0xc00 into the upper 16 bits of register r3  
0000076c 3c600c00 lis r3, 0xc00 {0xc000000}  
  
; Move to Special Purpose Register:  
; Move the contents of register r3 to special-purpose register 0x230  
00000770 7c708ba6 mtspr 0x230, r3  
  
; Instruction Synchronize:  
; Synchronize the instruction stream, ensuring all previous instructions  
    are completed before continuing  
00000774 4c00012c isync
```

```
; Load Immediate Shifted:  
; Load the immediate value 0xc00 into the upper 16 bits of register r3  
0000076c 3c600c00 lis r3, 0xc00 {0xc000000}  
  
; Move to Special Purpose Register:  
; Move the contents of register r3 to special-purpose register 0x230  
00000770  
works well for assembly snippets  
  
; Instruction Synchronize:  
; Synchronize the instruction stream, ensuring all previous instructions  
    are completed before continuing  
00000774 4c00012c isync
```

# Use Case Identification of Algorithms

“Analyze the following code and explain what it does.”

This function performs some kind of **data transformation** or **encryption/decryption** by manipulating bytes and **using loops to iterate through memory**.

This function performs some kind of data transformation or  
encryption/decryption through memory

often generic or wrong

oops to iterate

# Use Case Code Logic Explanation

“The following code implements a crypto algorithm. Which one and why?”

## RC4 KSA

The code snippet likely implements the **RC4 cryptographic algorithm** due to its characteristic **key scheduling** and pseudo-random generation processes.

## RC4 KSA

The code snippet likely implements the RC4 cryptographic algorithm due to its characteristic key context often helps generation processes.

# Use Case Verification Oracle

“Does the following code implement XXX? Answer YES or NO.”

*"Does the following code implement VVV? Answer YES or NO"*

good to filter false positives from heuristics

# Use Case Script Development

“The following code implements a string decryption routine.  
Rebuilt it in Python.”

```
def decrypt(input_string):
    xor_key = 0x22
    decrypted_chars = [chr(ord(char) ^ xor_key) for char in input_string]
    return ''.join(decrypted_chars)
```

```
def decrypt(input_string):
    xor_key = 0x22
    decrypted_chars = [chr(ord(char) ^ xor_key) for char in input_string]
    return ''.join(decrypted_chars)
```

works fine for small tasks

- good for **approximations** and **high-level** understanding
- answers sometimes **generic** or **wrong**
- **adding context** often helps
- **no** real (code) understanding

# Tools

## Tools & Integrations (Selection)

- various tools and wrappers for **commercial LLMs**
- **IDA Pro**: Gepetto, Copilot for IDA Pro
- **Ghidra**: GhidraChatGPT, GptHidra
- **Binary Ninja**: Sidekick, BinaryNinja-OpenAI

# Gepetto

The screenshot shows the IDA Pro interface with several windows open:

- [2] IDA View-A
- [3] Pseudocode-A
- [4] Hex View-1
- [5] Structures
- [6] Enums

The assembly view (View-A) contains the following code:

```
1 int sub_100011AE()
2 {
3     int result; // eax
4     char Str[1021]; // [esp+4h] [ebp-504h] BYREF
5     __int16 v2; // [esp+401h] [ebp-107h]
6     char v3; // [esp+403h] [ebp-105h]
7     char Destination[257]; // [esp+404h] [ebp-104h] BYREF
8     __int16 v5; // [esp+505h] [ebp-3h]
9     char v6; // [esp+507h] [ebp-1h]
10
11     memset(Destination, 0, sizeof(Dest
12     v5 = 0;
13     v6 = 0;
14     strcat(Destination, PathName);
15     strcat(Destination, aUpgradeExe);
16     result = access(Destination, 0);
17     if ( result == -1 )
18     {
19         memset(Str, 0, sizeof(Str));
20         v2 = 0;
21         v3 = 0;
22         strcat(Str, aBgkkhwr0dr1d80);
```

A context menu is open at the end of the function body, specifically over the closing brace of the function definition. The menu items are:

- Add breakpoint F2
- Synchronize with
- Edit Ivar comment... /
- Collapse declarations Numpad++
- Mark as decompiled
- Copy to assembly
- Hide casts \
- Gepetto
  - Explain function Ctrl+Alt+G
  - Rename variables Ctrl+Alt+R
- Font...

<https://github.com/JusticeRage/Gepetto>

# Binary Ninja Sidekick

The screenshot shows the Binary Ninja interface with the following details:

- Project:** minia\_arm
- File Type:** ELF - Linear - High Level IL
- Function:** void\* sub\_10778(char\* arg1, char\* arg2, int32\_t arg3)
- Assembly View:** The assembly code for the function is displayed, showing a loop that XORs each byte of the input strings (arg1 and arg2) with arg3. The code uses registers r0.3, r0.5, and r0.7, and memory locations data\_lebfc and data\_lec28.
- Python View:** A Python code editor on the left shows a simplified Python function for decrypting strings. It takes an input string, initializes a decrypted string, and loops through each character, XORing it with arg3 and appending it to the decrypted string.
- Suggestions:** A sidebar lists various pointers and symbols such as arg1, r0.1, r0.3, r0.5, r0.7, arg2, arg3, r2.3, r2.5, r2.6, r2.9, r2.10, r2.16, r7.1, r11, lr, and dataBufferIndex.
- Strings:** A search bar at the bottom allows searching for specific strings within the assembly code.

<https://sidekick.binary.ninja/>

# Downsides



Internet connection required



Every query costs \$\$\$



Privacy risks



## Local LLMs

### Pros:

- offline
- privacy-sensitive

### Pros:

- offline
- privacy-sensitive

### Cons:

- slower

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### Cons:

- slower
- less powerful

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- offline
- privacy-sensitive

### Cons:

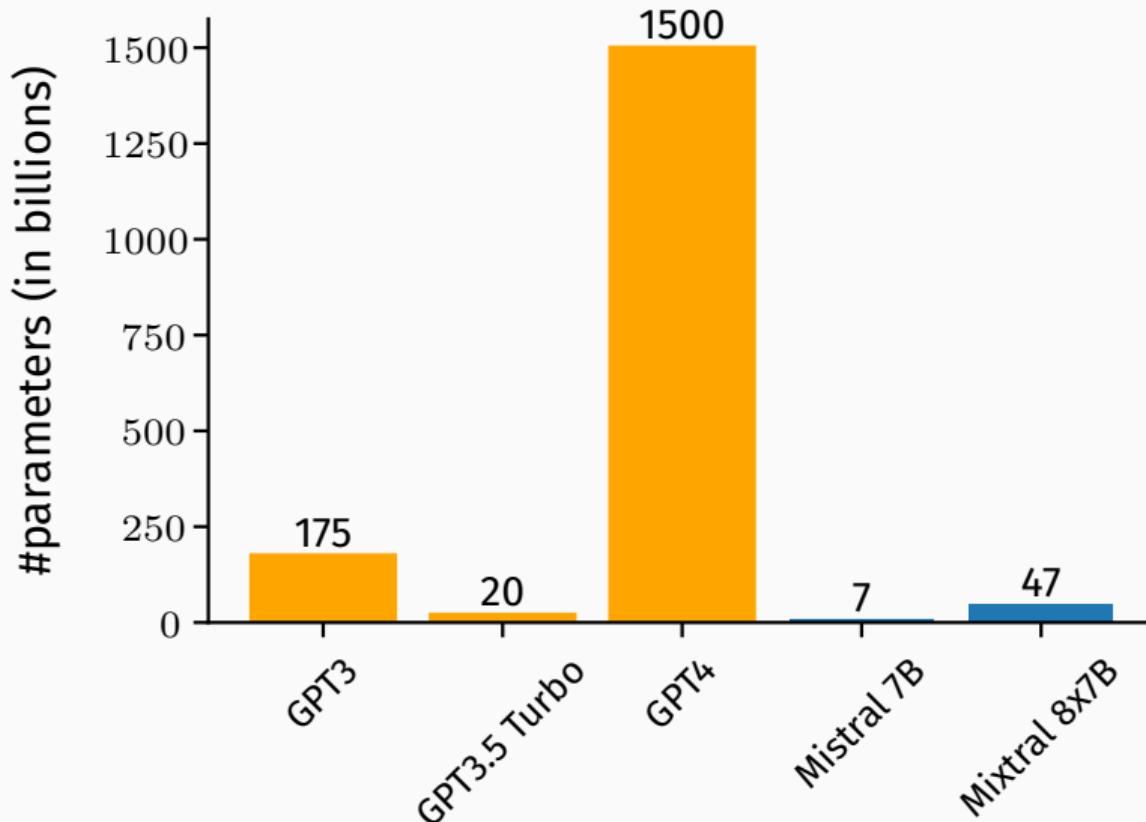
- slower
- less powerful
- computation resources

Local LLMs are .. slower

Anecdotally, for some function renaming queries:

- GPT query: <2s
- Mistral 7B on M1 Macbook Pro: 5-8s
- Mistral 7B on M3 Macbook Air: 10s

Local LLMs are .. less powerful



### Good GPU:

- NVIDIA GTX 3090
- NVIDIA GTX 4090

### ARM-based Mac:

- M1/M2/M3 Macbook
- or workstation

Good news:  
We can already use local LLMs for RE

## (Unfair) Comparison to GPT4

Use Case	Mistral 7B	Mixtral 8x7B
function renaming	✓	✓
identify library functions	✗	✓
enhance decompiler output	✗	✓
annotate assembly	✗	✓
explain code logic	✗	✓
helper script development	✗	✓
verification oracle	✗	✓

Can we do better?

## Context-sensitive Annotations

Context helps..

..so let's use available information!

## Context-sensitive Annotations

⇒ incorporate insights from static analysis

example: function renaming



rename only “relevant” functions

## Context-sensitive Annotations



improves accuracy and speed

# Context-sensitive Annotations

```
1000ead0 enum WIN32_ERROR __fastcall configure_proxy_settings(void* arg1)

1000ead0    int32_t var_4_2 = 0xffffffff
1000ead2    int32_t (* var_8)() = sub_1003bf4b
1000eadd    int32_t* fsbase
1000eadd    int32_t var_c = *fsbase
1000ea0e    *fsbase = &var_c
1000eaee    int32_t var_114 = 0
1000eb08    int32_t var_4 = 0
1000eb1a    if (sub_10005d00(&var_114, 0x80000001, "Software\Microsoft\Windows\CurrentVersion\Internet Settings\ProxyServer", 1, 0) != 0)
1000eb2f        int32_t var_110
1000eb2f        sub_10005db0(&var_114, "ProxyEnable", &var_110)
1000eb3a        if (var_110 != 0)
1000eb4a            sub_10005db0(&var_114, "ProxyServer", &var_110)
1000eb59            char* eax_3 = sub_1002e3d0(&var_110, "http=")
1000eb65            void* esi_2
1000eb65            if (eax_3 != 0)
1000eb6d                esi_2 = &eax_3[5]
```

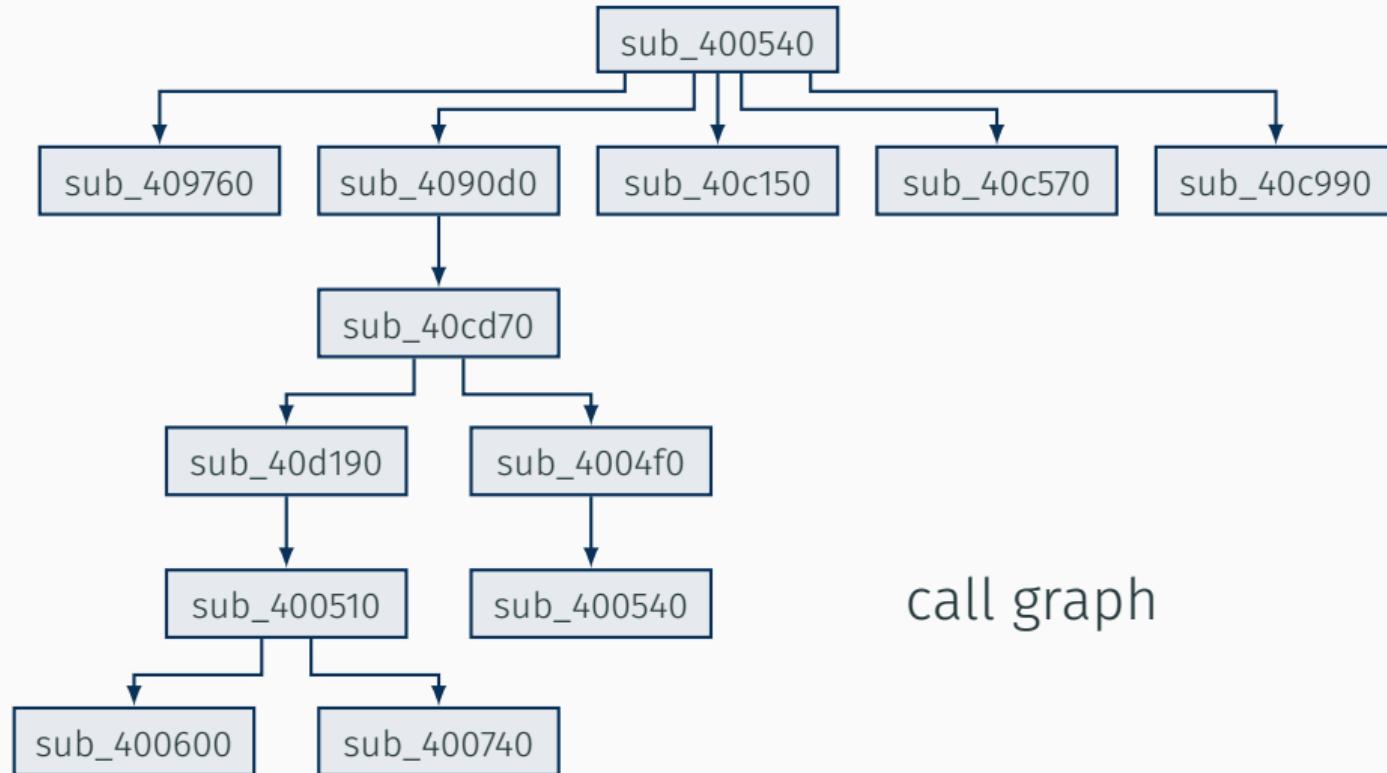
# Context-sensitive Annotations

```
1000ead0 enum WIN32_ERROR __fastcall configure_proxy_settings(void* arg1)

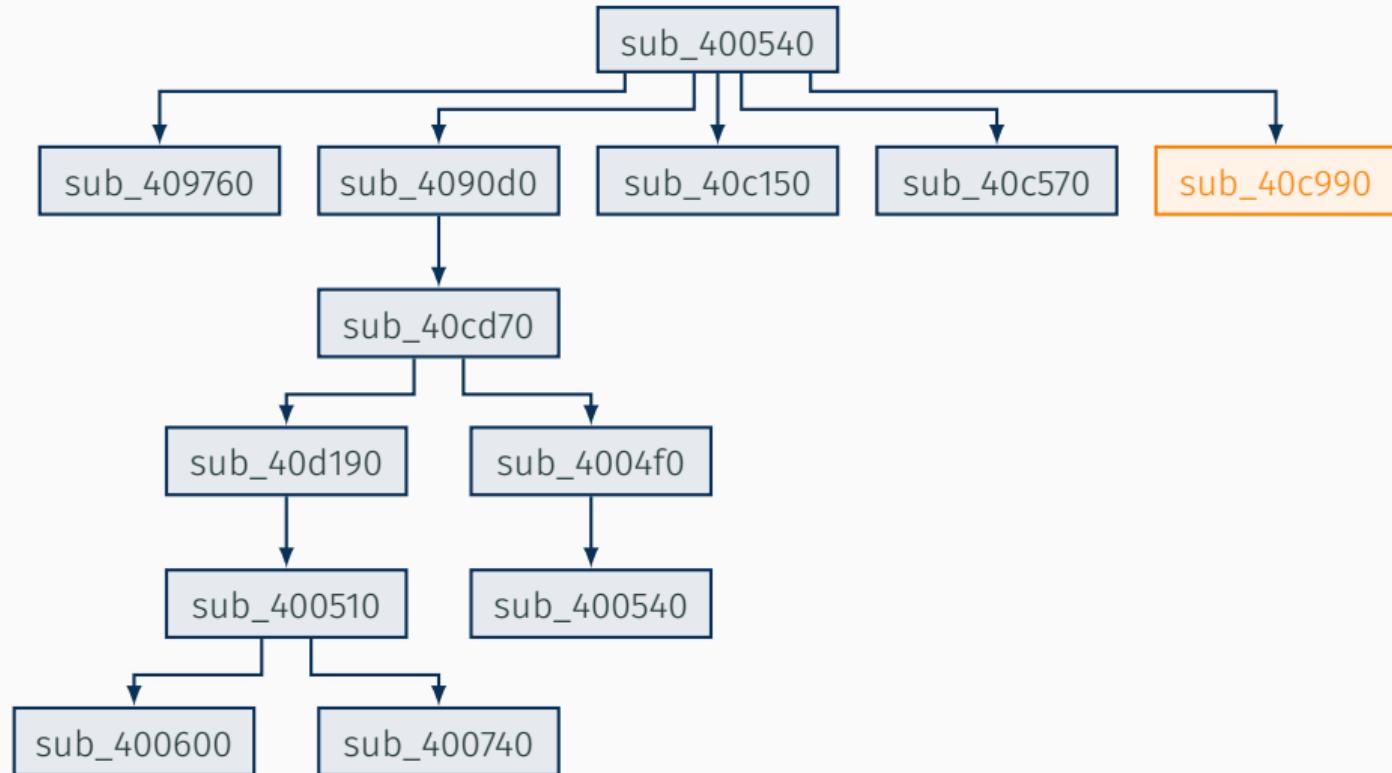
1000ead0    int32_t var_4_2 = 0xffffffff
1000ead2    int32_t (* var_8)() = sub_1003bf4b
1000eadd    int32_t* fsbase
1000eadd    int32_t var_c = *fsbase
1000eaade   *fsbas
1000eaee    int32_
1000eb08    int32_
1000eb1a    if (sub_10005d00(&var_114, 0x80000001, "Software\Microsoft\Windows\CurrentVersion\Internet Settings\ProxyServer", "ProxyEnable", &var_110) != 0)
1000eb2f    int32_t var_110
1000eb2f    sub_10005db0(&var_114, "ProxyEnable", &var_110)
1000eb3a    if (var_110 != 0)
1000eb4a    sub_10005db0(&var_114, "ProxyServer", &var_110)
1000eb59    char* eax_3 = sub_1002e3d0(&var_110, "http=")
1000eb65    void* esi_2
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1000eb6d    |    esi_2 = &eax_3[5]
```

strings and API functions

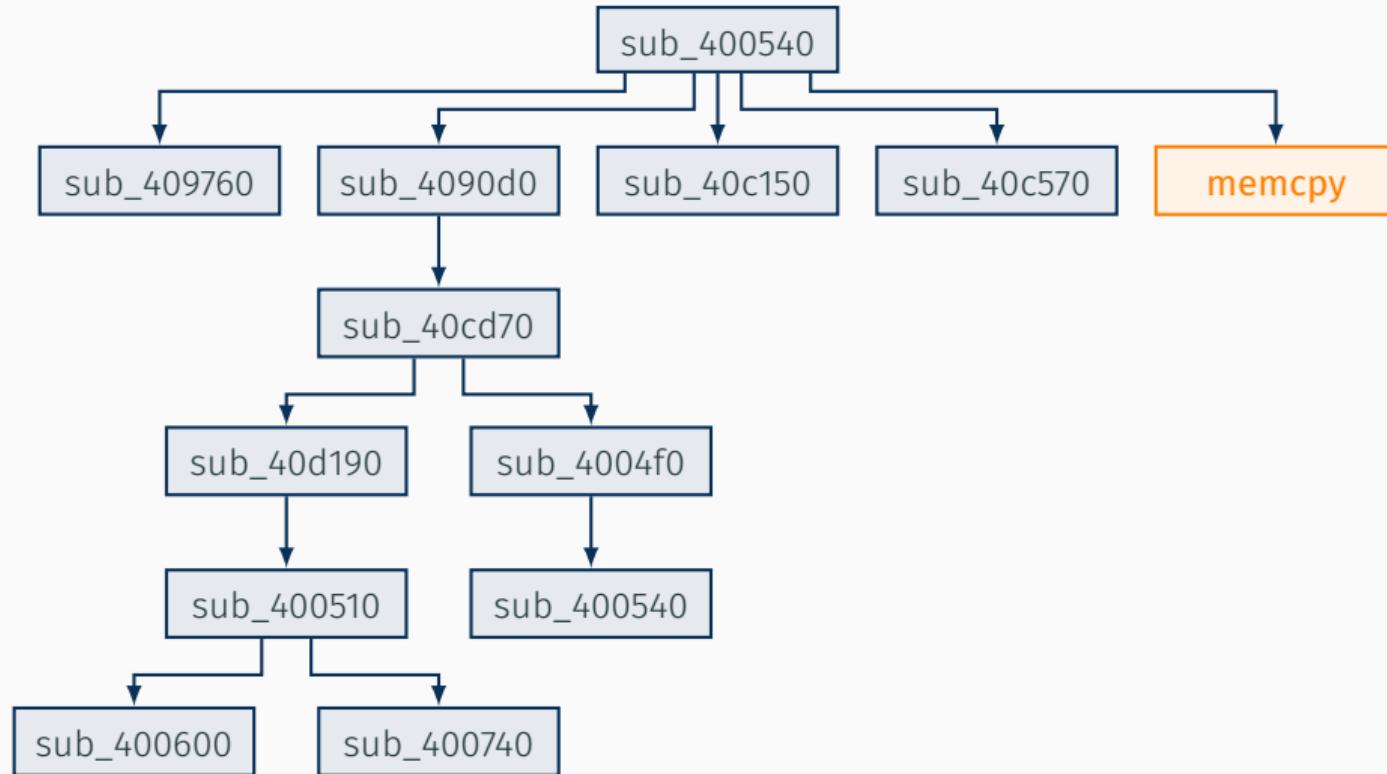
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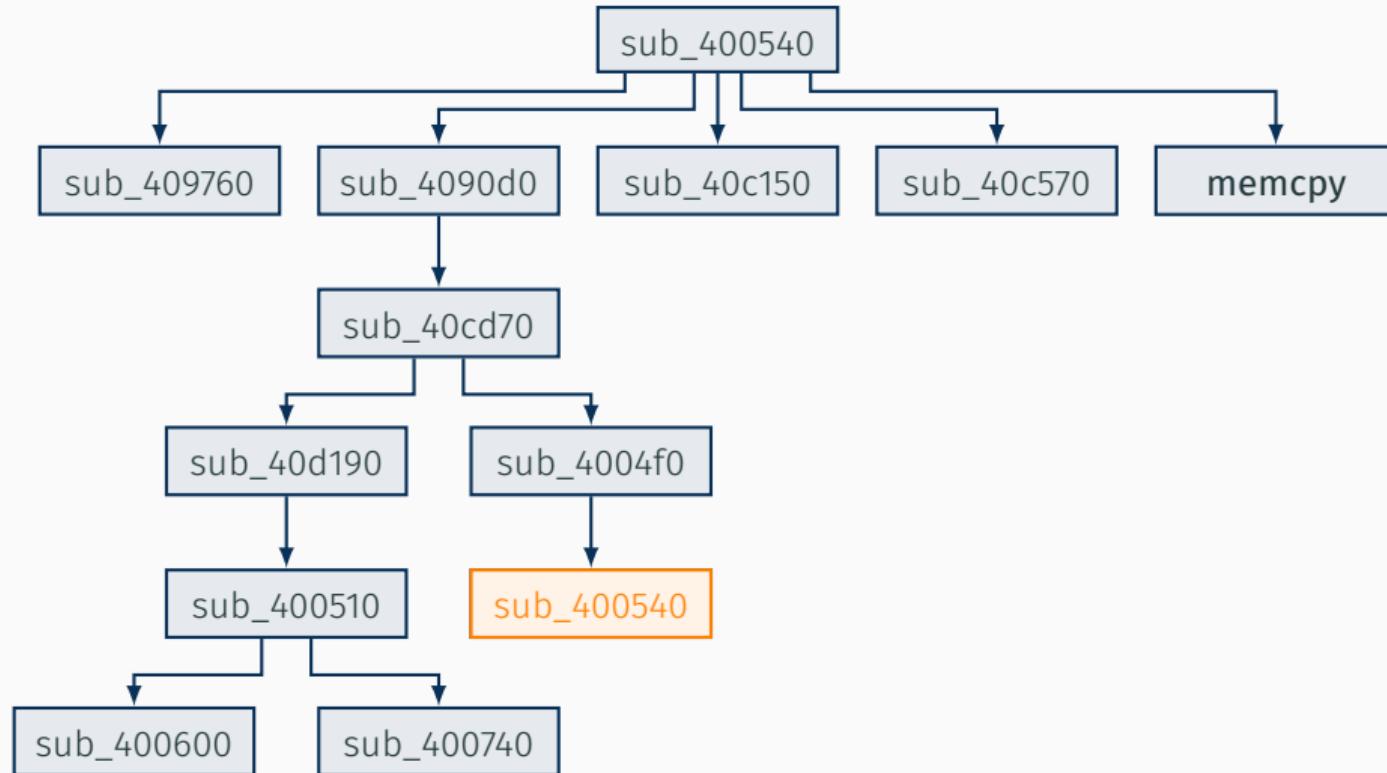
## Context-sensitive Annotations



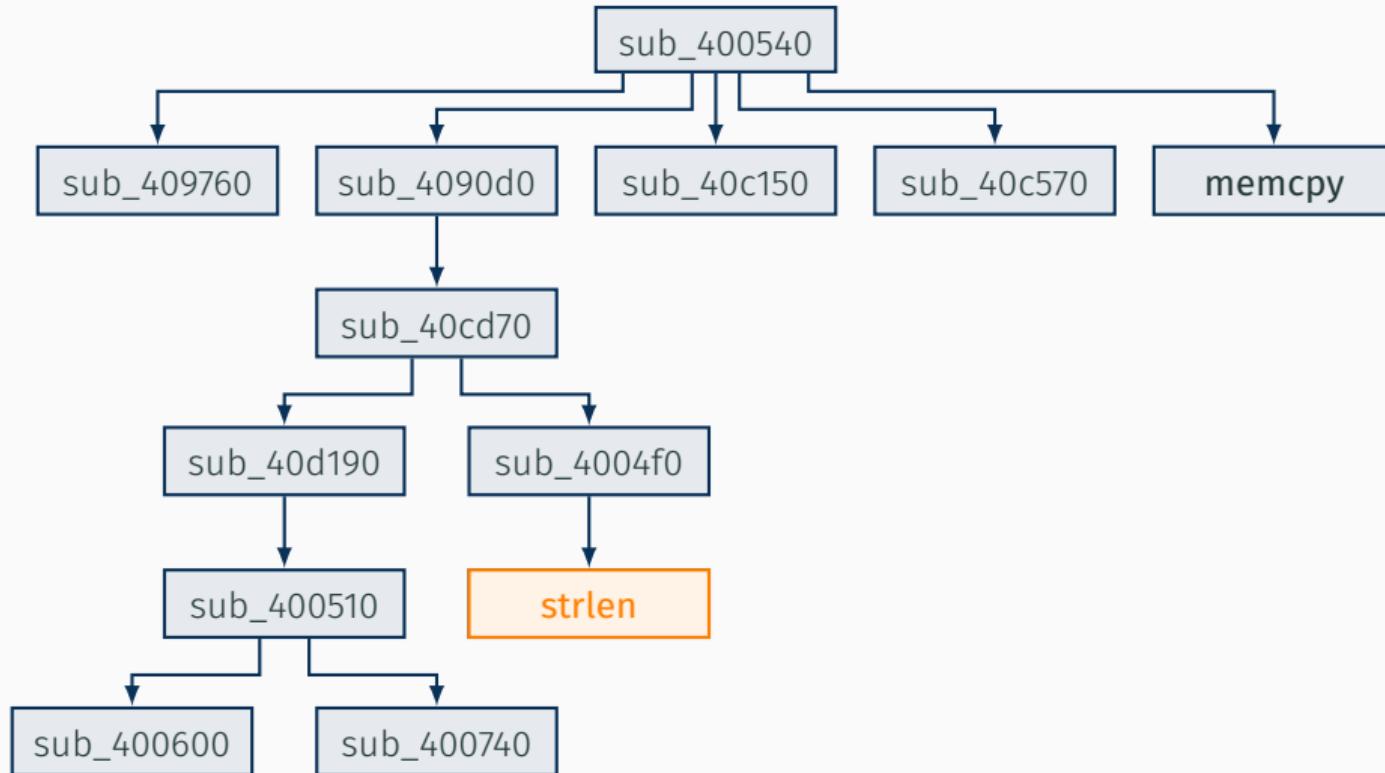
## Context-sensitive Annotations



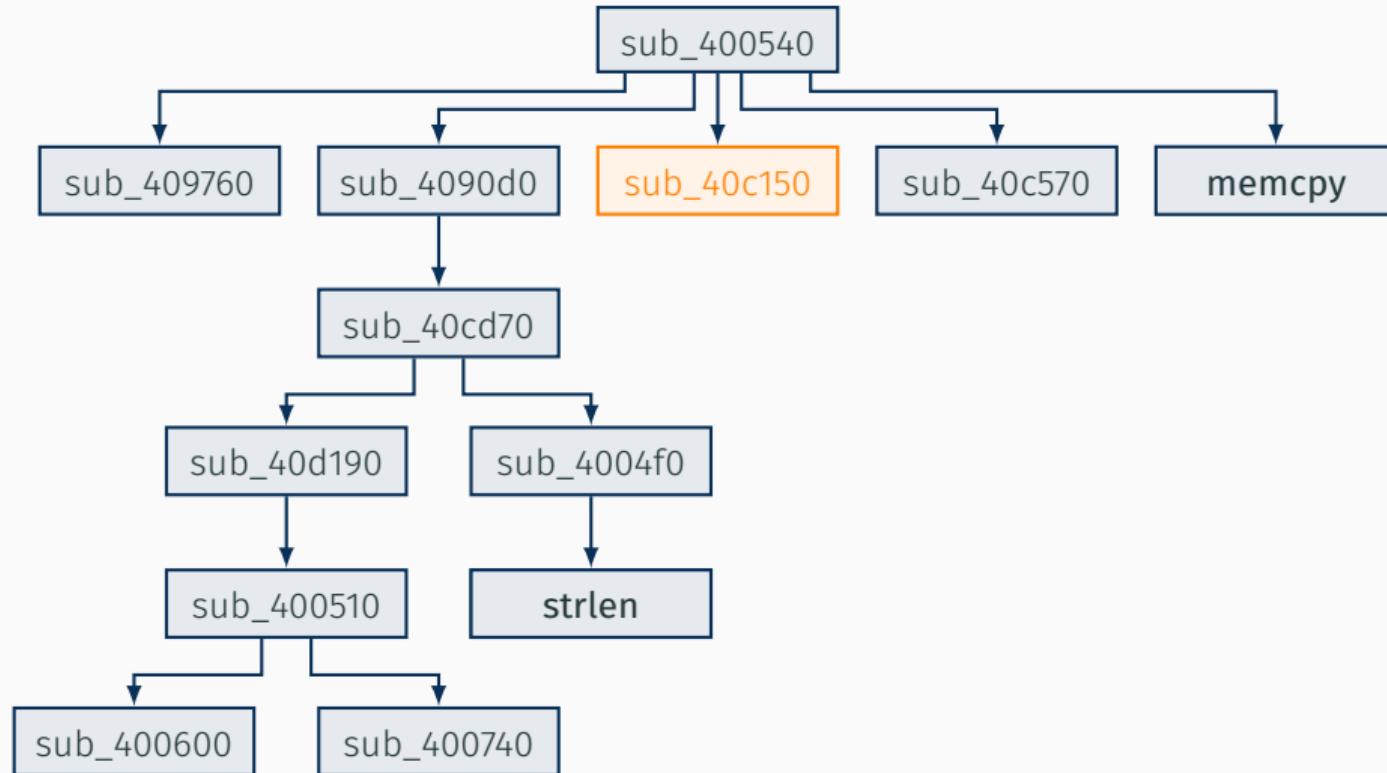
## Context-sensitive Annotations



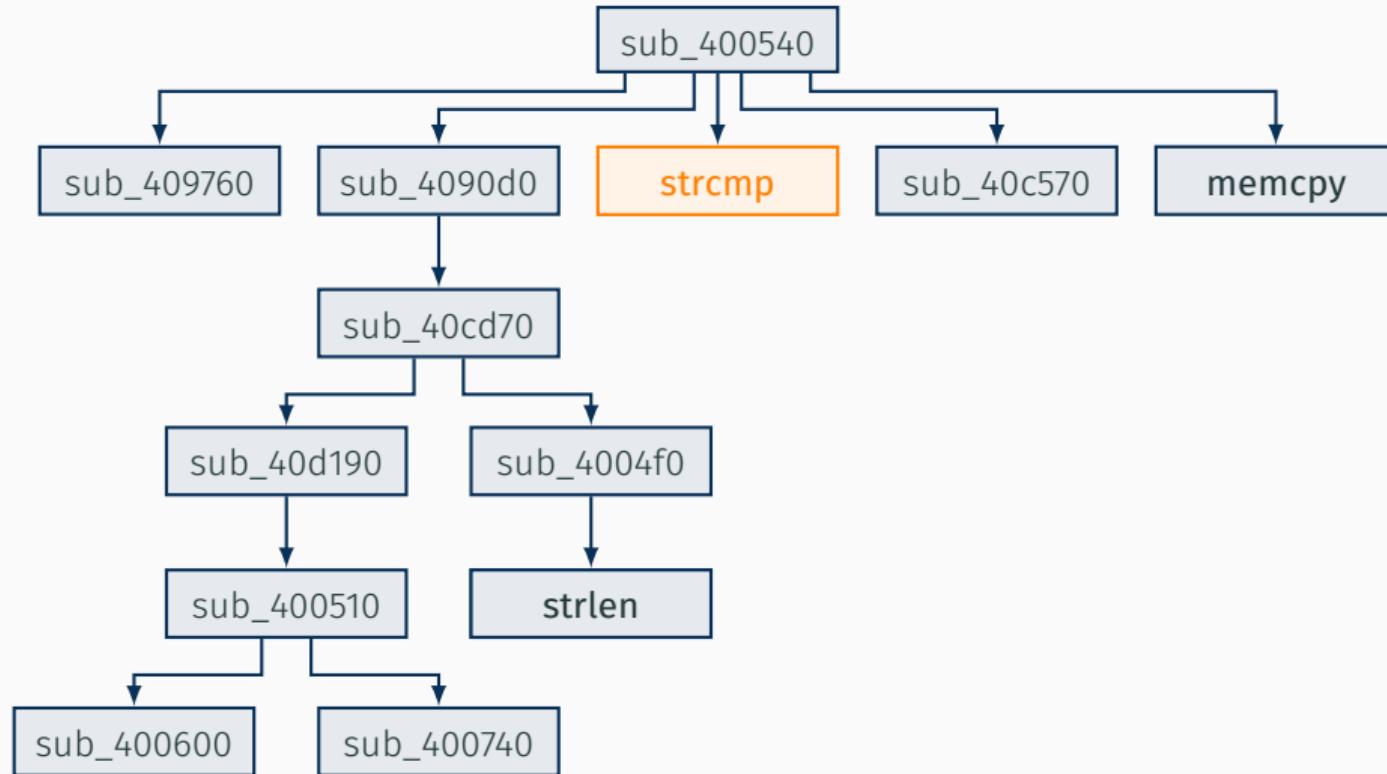
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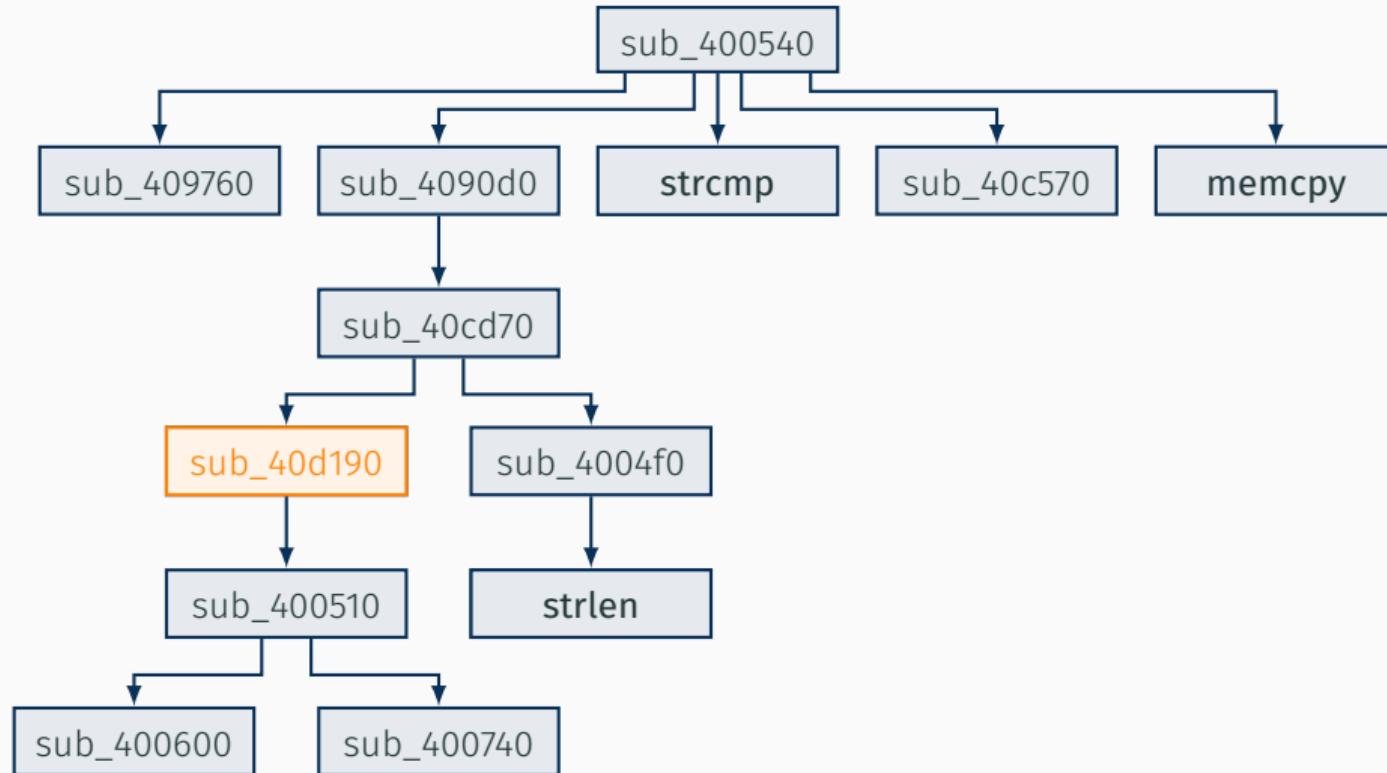
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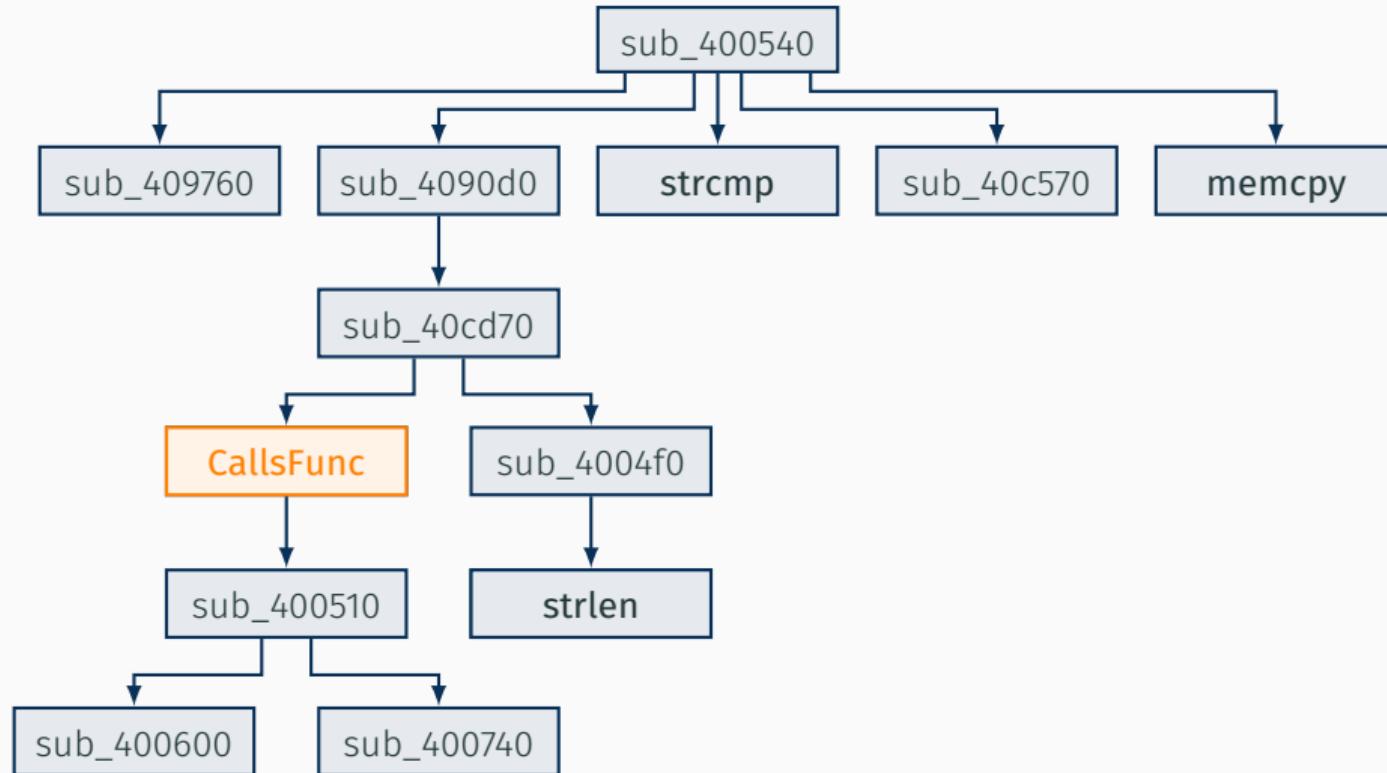
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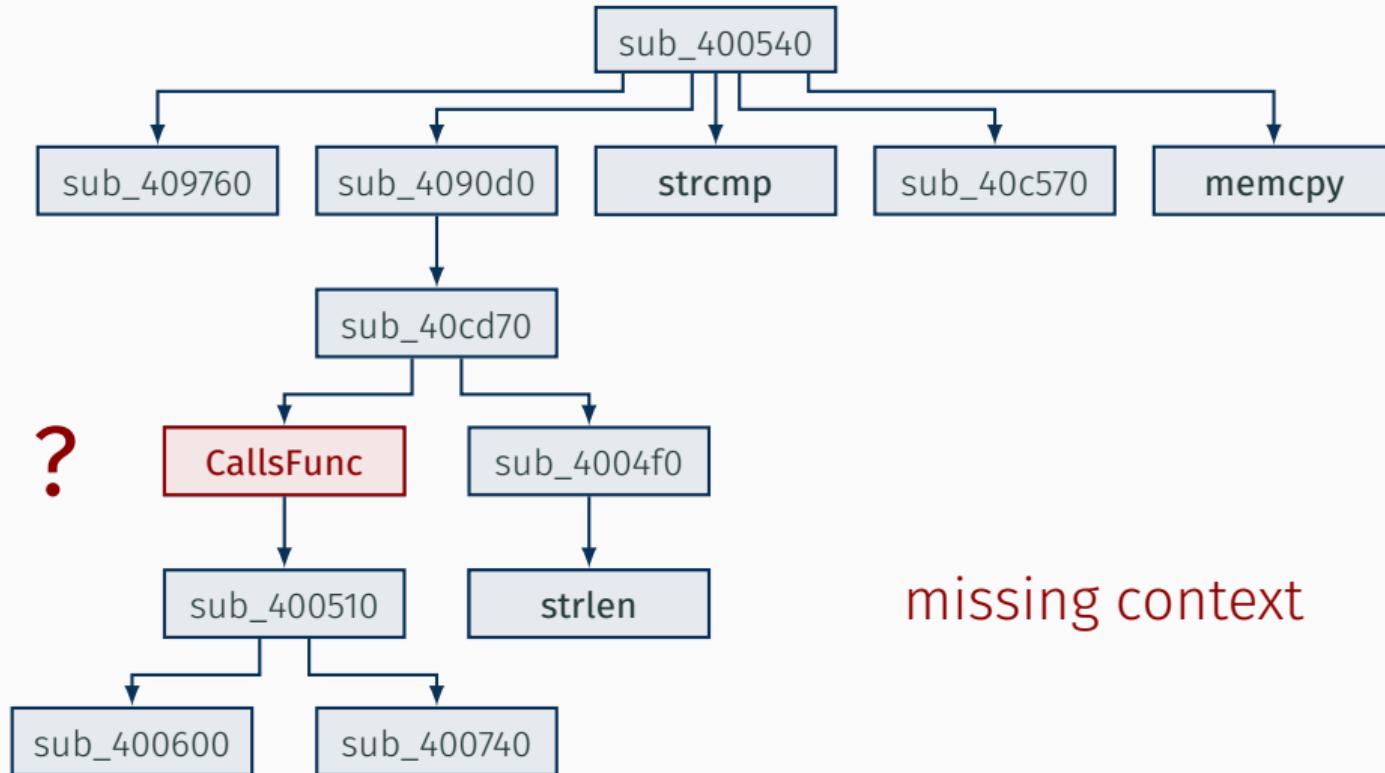
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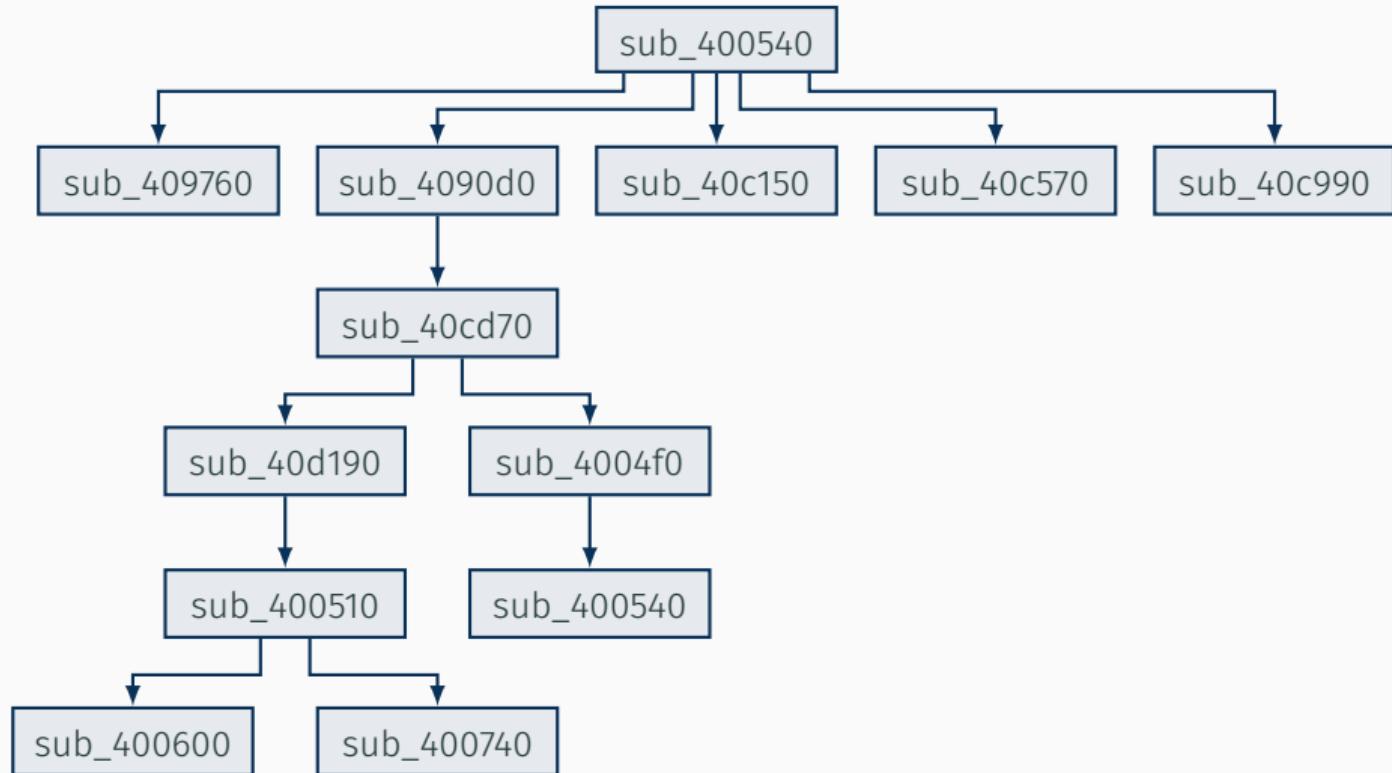
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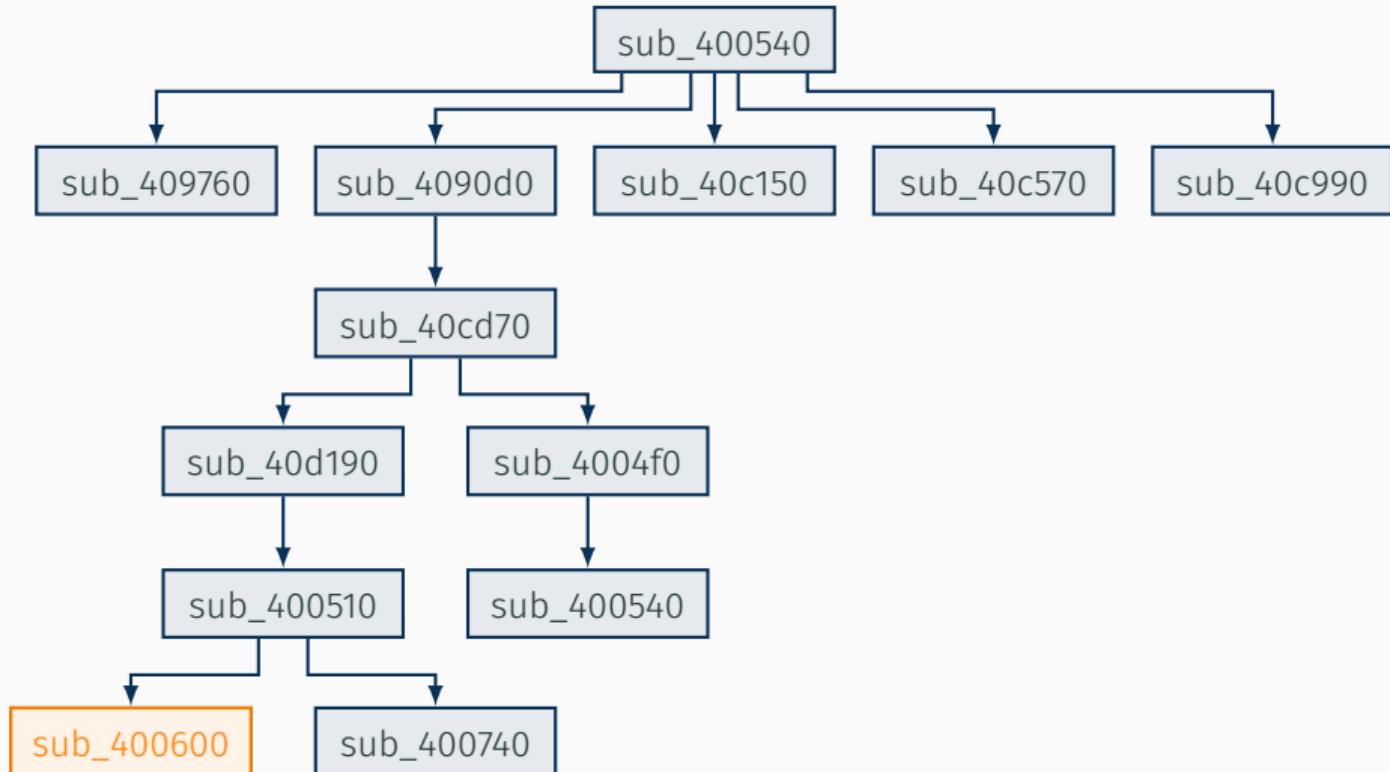
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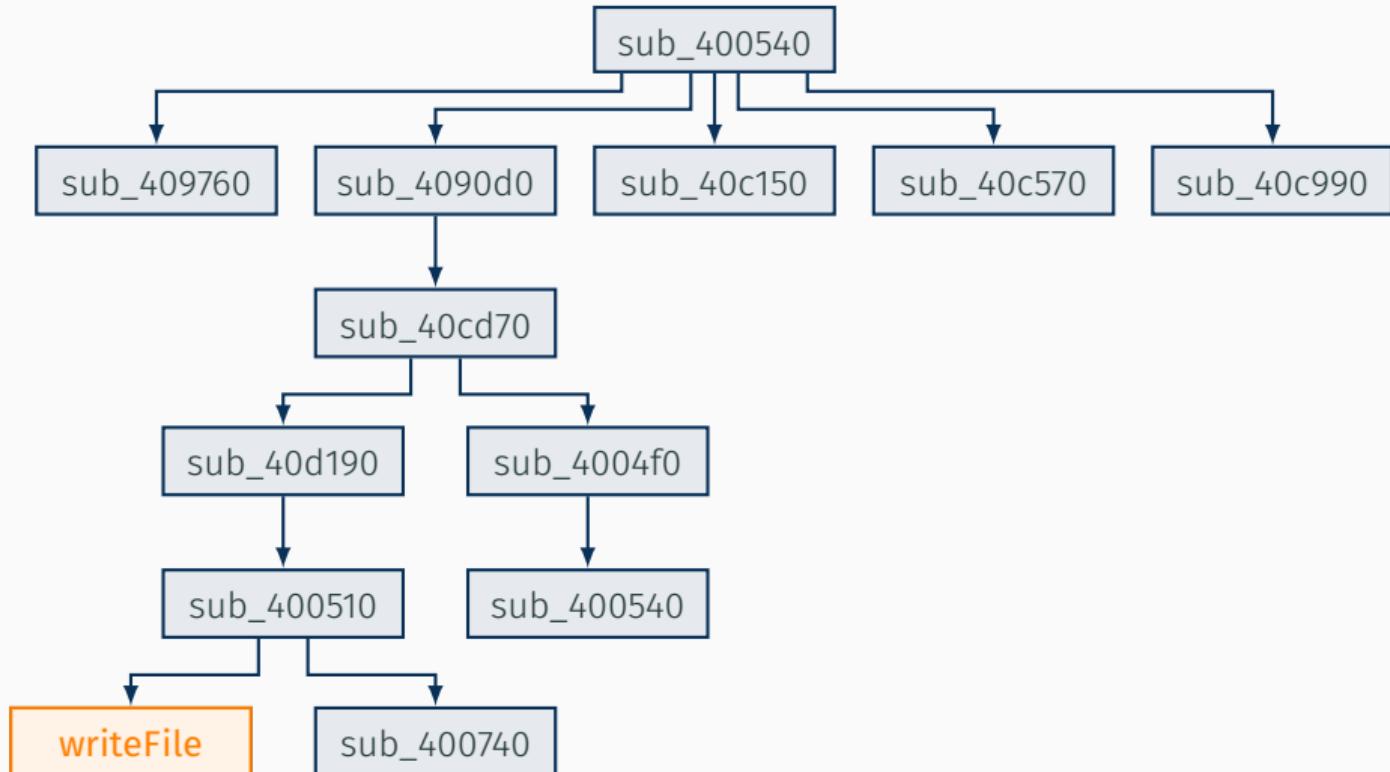
## Context-sensitive Annotations: Bottom-up Propagation



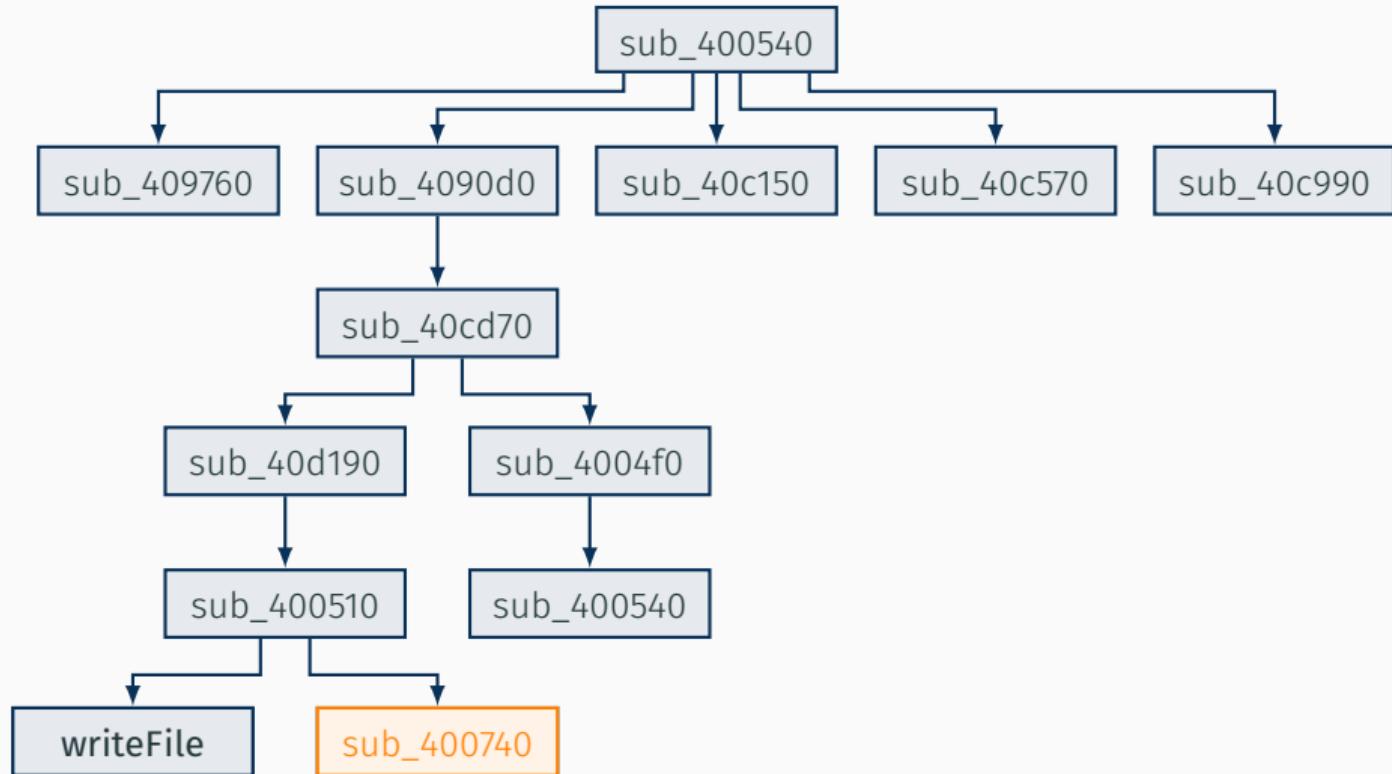
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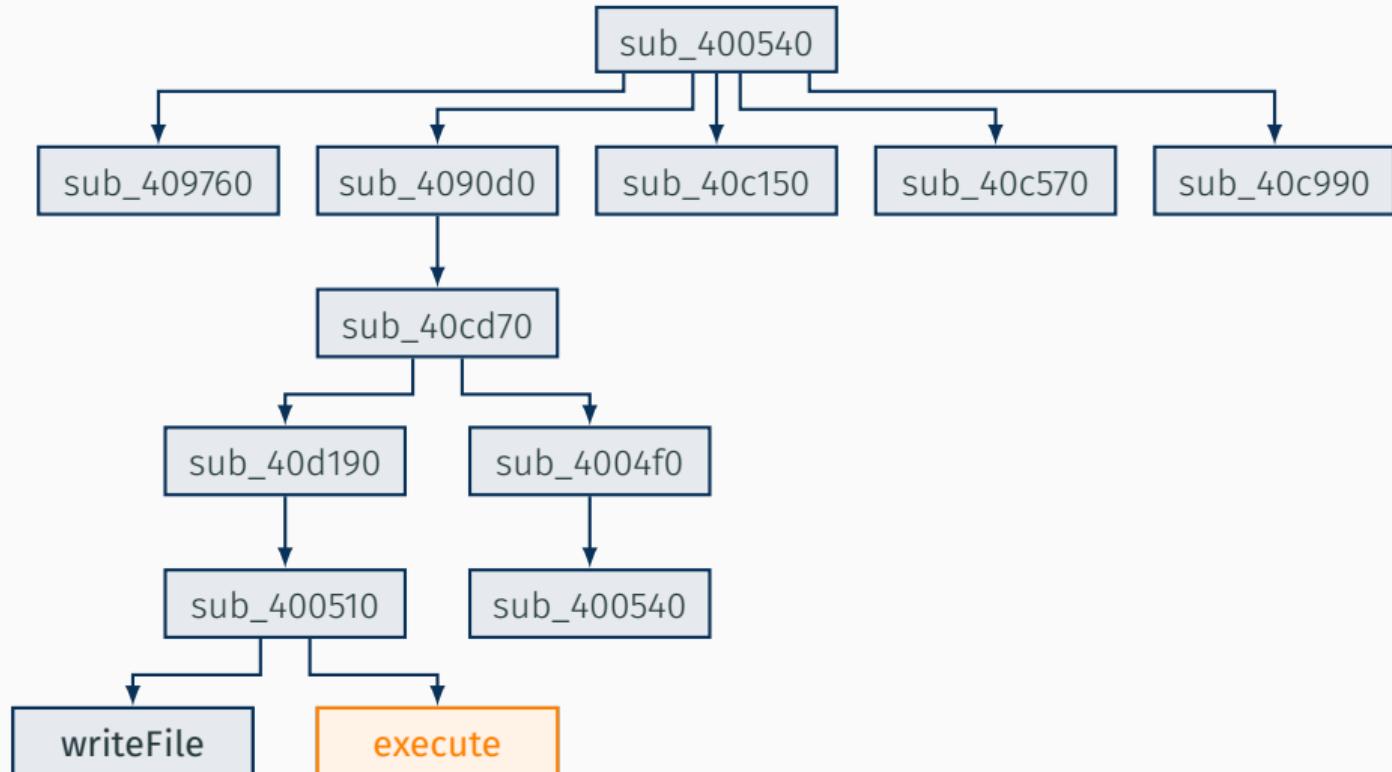
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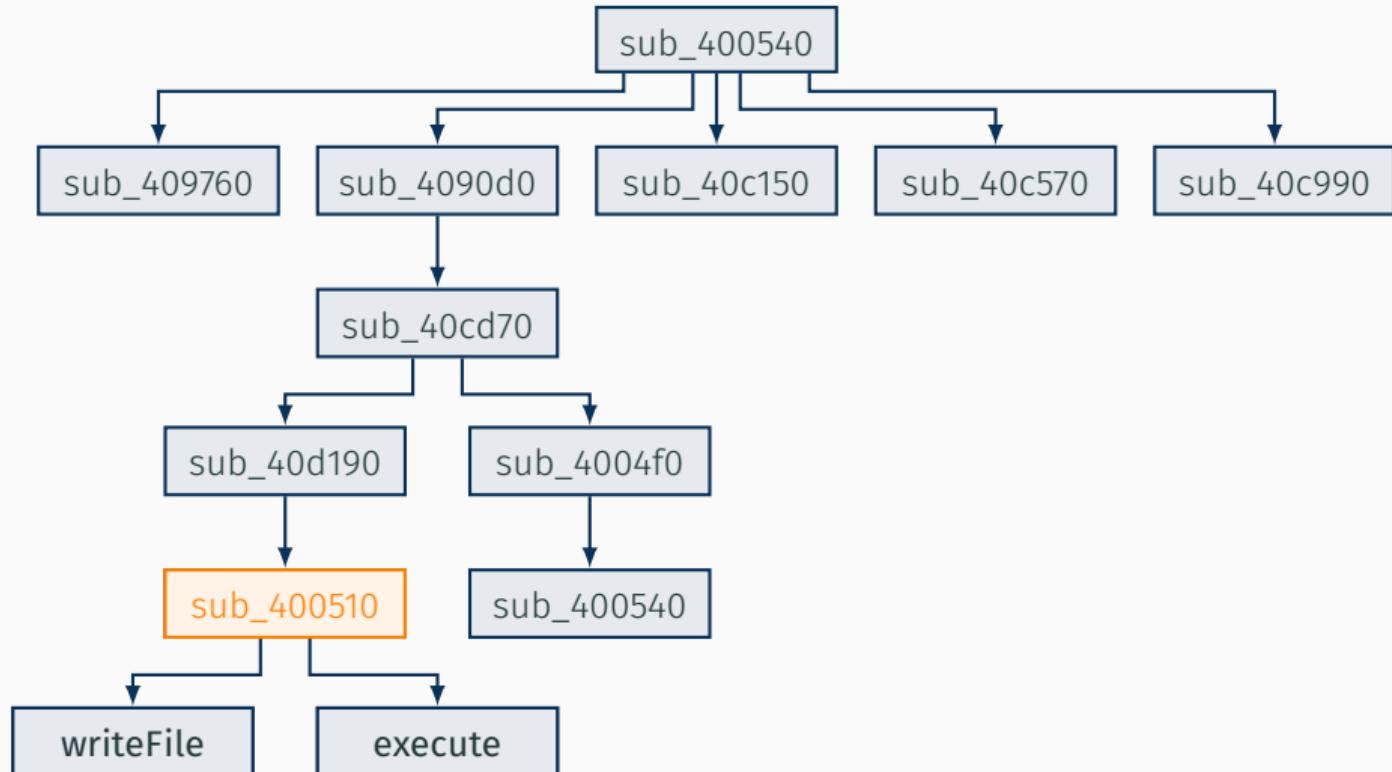
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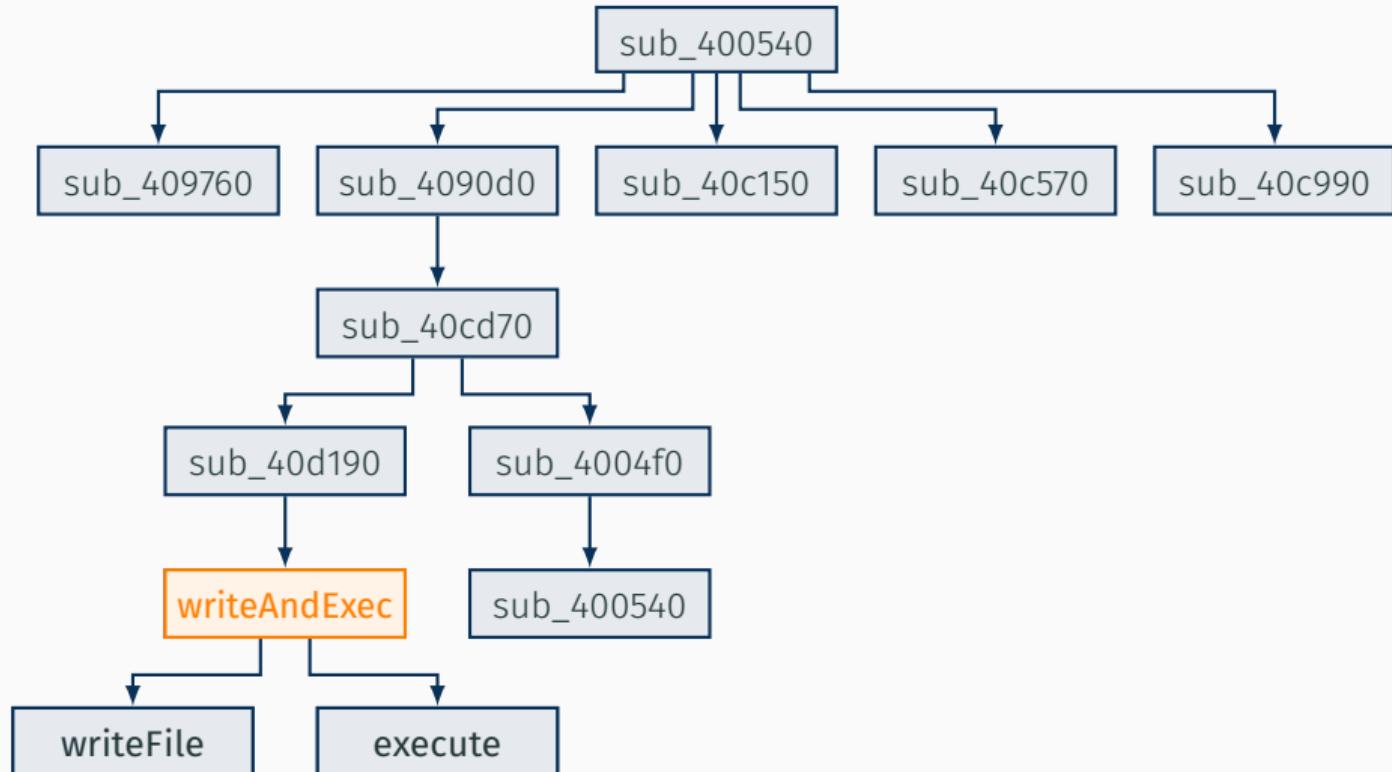
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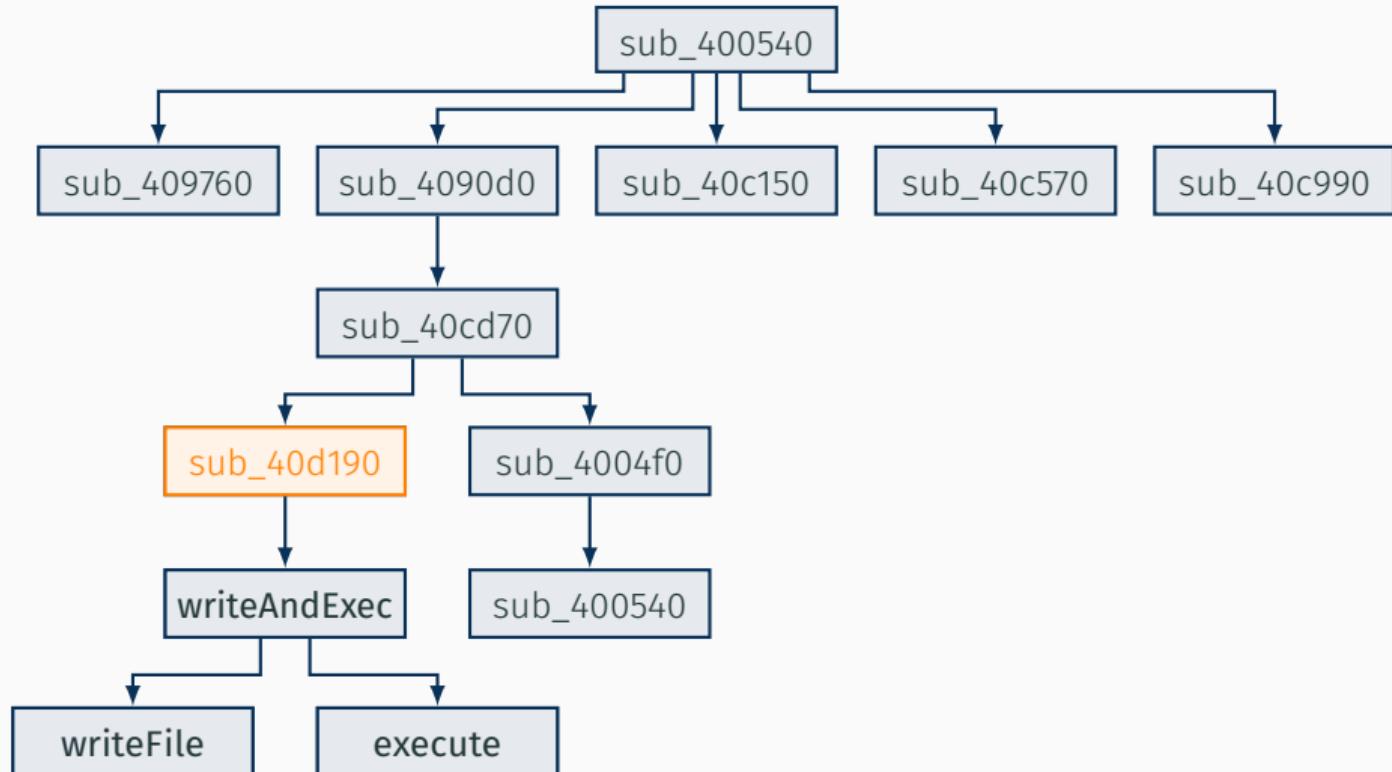
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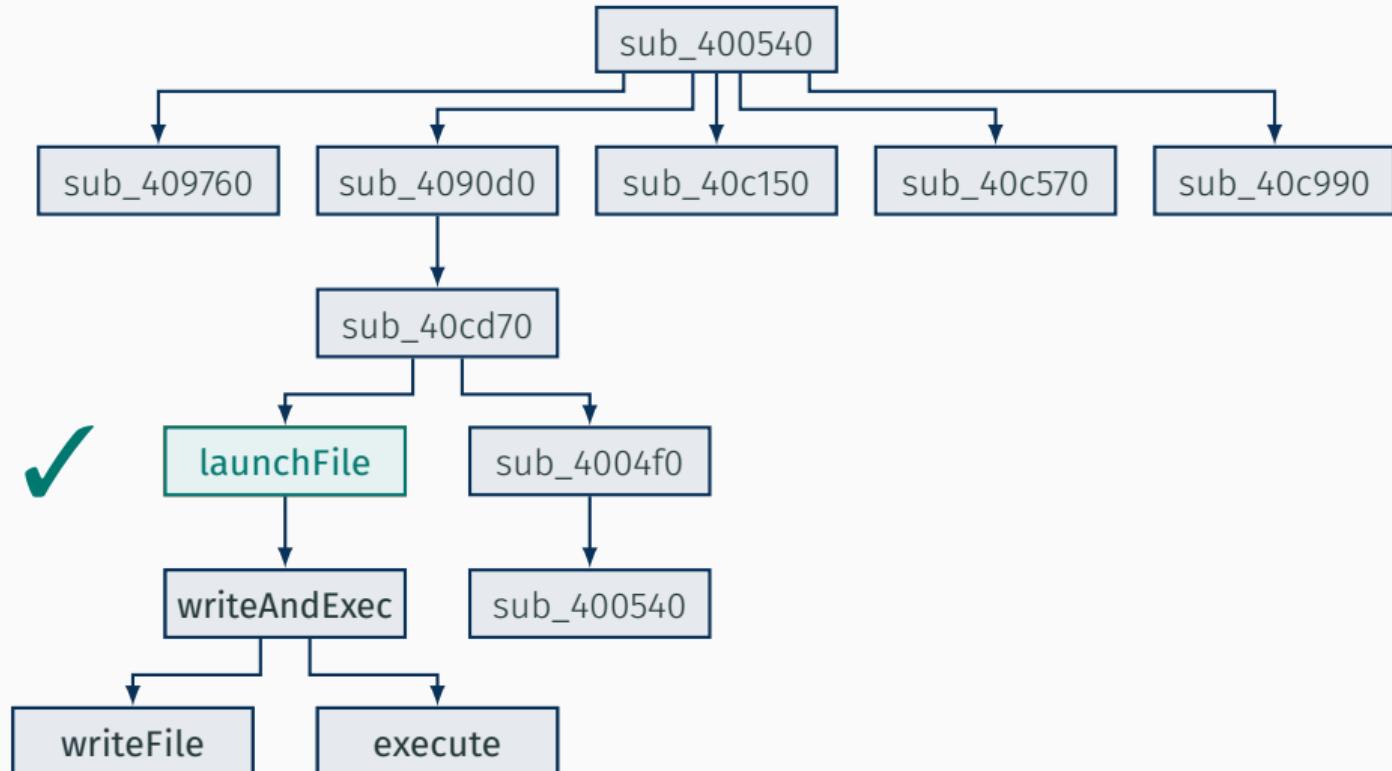
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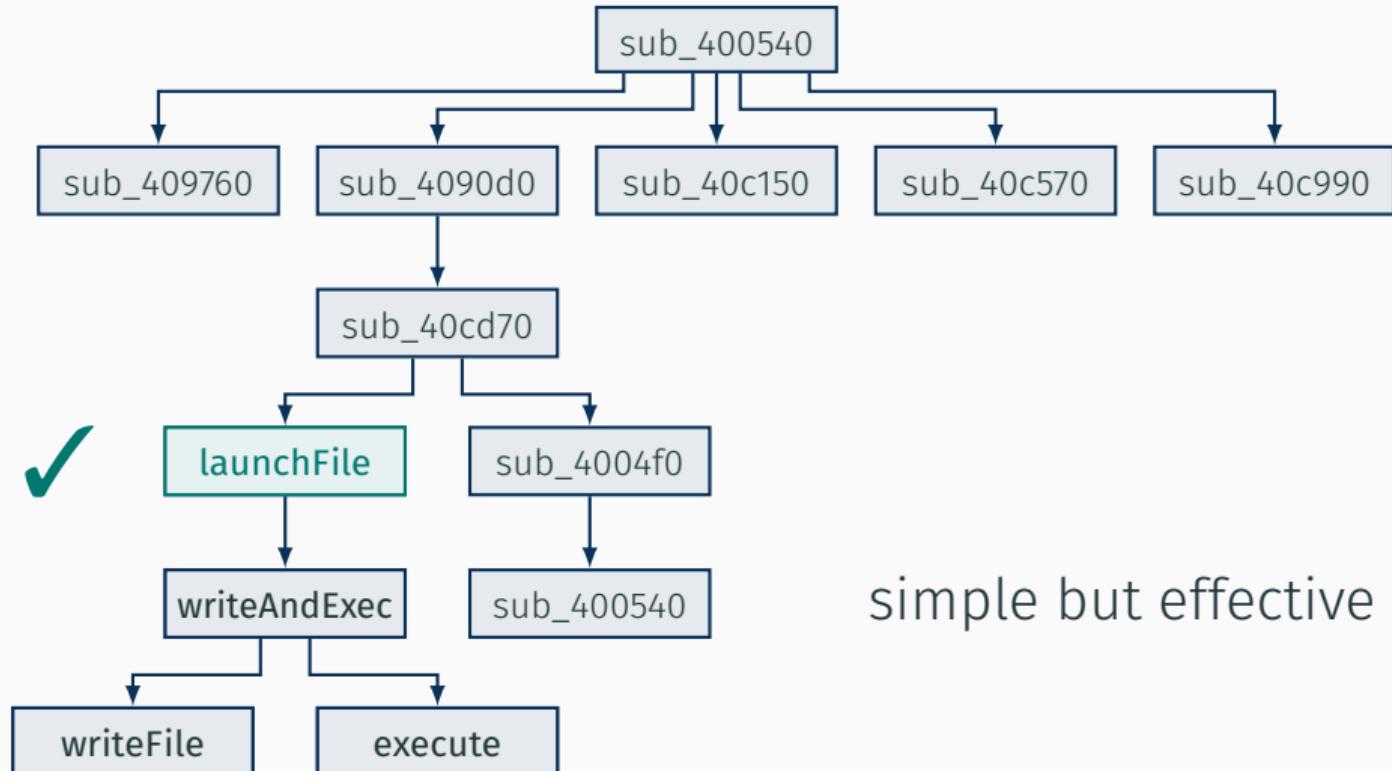
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## Context-sensitive Annotations: Bottom-up Propagation



# Tools

The screenshot shows a GitHub repository page for 'ReverserAI (v1.1)'. At the top, there are links for 'README' and 'GPL-2.0 license'. To the right are icons for edit, fork, and more. The main title 'ReverserAI (v1.1)' is displayed in large bold letters. Below it, the author is listed as 'Author: Tim Blazytko'. A descriptive text follows: 'Provides automated reverse engineering assistance through the use of local large language models (LLMs) on consumer hardware.' A section titled 'Description:' is present, containing a detailed explanation of the project's purpose and architecture.

**ReverserAI (v1.1)**

Author: Tim Blazytko

*Provides automated reverse engineering assistance through the use of local large language models (LLMs) on consumer hardware.*

**Description:**

ReverserAI is a research project designed to automate and enhance reverse engineering tasks through the use of locally-hosted large language models (LLMs). Operating entirely offline, this initial release features the automatic suggestion of high-level, semantically meaningful function names derived from decompiler output. ReverserAI is provided as a Binary Ninja plugin; however, its architecture is designed to be extended to other reverse engineering platforms such as IDA and Ghidra.

[https://github.com/mrphrazer/reverser\\_ai](https://github.com/mrphrazer/reverser_ai)

- Ninja plugin to include local LLMs
- more playground than finished product
- supports two models (Mistral 7B and Mixtral 8x7B)

Can we do better?

## Things to improve

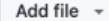
- better queries (prompt engineering)
- better models: <https://github.com/eugeneyan/open-llms>
- fine-tuned models

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 **aidapal** Public

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 main  1 Branch  0 Tags  Go to file  Add file  Code

**AverageBusinessUser** Update idapal\_qt\_interface.py  274aa64 · yesterday  5 Commits

 README.md Update README.md 3 weeks ago

 idapal.py initial plugin upload 3 weeks ago

 idapal\_qt\_interface.py Update idapal\_qt\_interface.py yesterday

 README  

**aiDAPal**

About  
aiDAPal is an IDA Pro plugin that uses a locally running LLM that has been fine-tuned for Hex-Rays pseudocode to assist with code analysis.

 Readme  
 Activity  
 Custom properties  
 94 stars  
 7 watching  
 9 forks  
Report repository

<https://github.com/atredispartners/aidapal>

- plugin for IDA with focus on enhancing decompiler output
- fine-tuned Mistral 7B ⇒ can keep up with GPT4
- also uses available context information

- plugin for IDA with focus on enhancing decompiler output
- fine-tuned **fine-tuning works well**<sup>PT4</sup>
- also uses available context information

## Currently Impossible

- *real* code analysis
- bug finding (beyond easy patterns)
- cross-function analysis

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- *real* code analysis
- bug finding (beyond easy patterns)  
**LLMs as helper, not automated analysts**
- cross-function analysis

# Future Trends

## Expectations

- enhanced scalability for broader inputs
- advancements in **on-device LLMs**

## Potential Emerging RE Applications

- semantic code search
- identification of noteworthy code segments
- patch recommendation systems
- binary similarity and clustering

# Conclusion

## Takeaways

1. LLMs are good for **approximations** and **high-level understanding**
2. But: They can be wrong and have **no** real (code) **understanding**
3. **Adding context** often helps increase accuracy
4. Local LLMs are somewhat worse but ensure **privacy**

## Summary

- (local) LLMs help, check them out
- but they are no panacea, be wary of the hype

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